

Learning quantum states and unitaries of bounded gate complexity

arXiv:2310.19882

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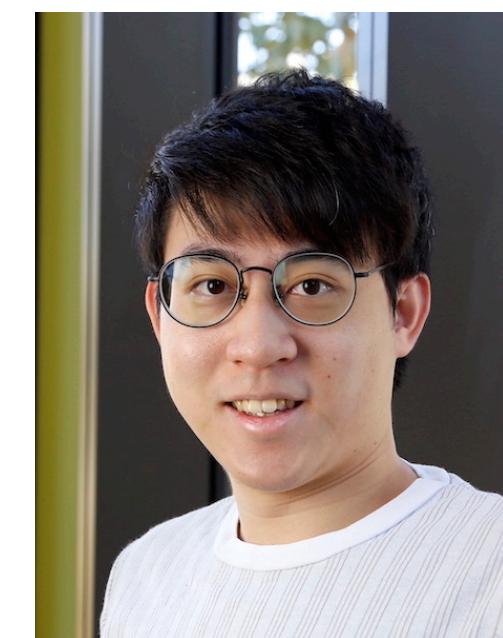
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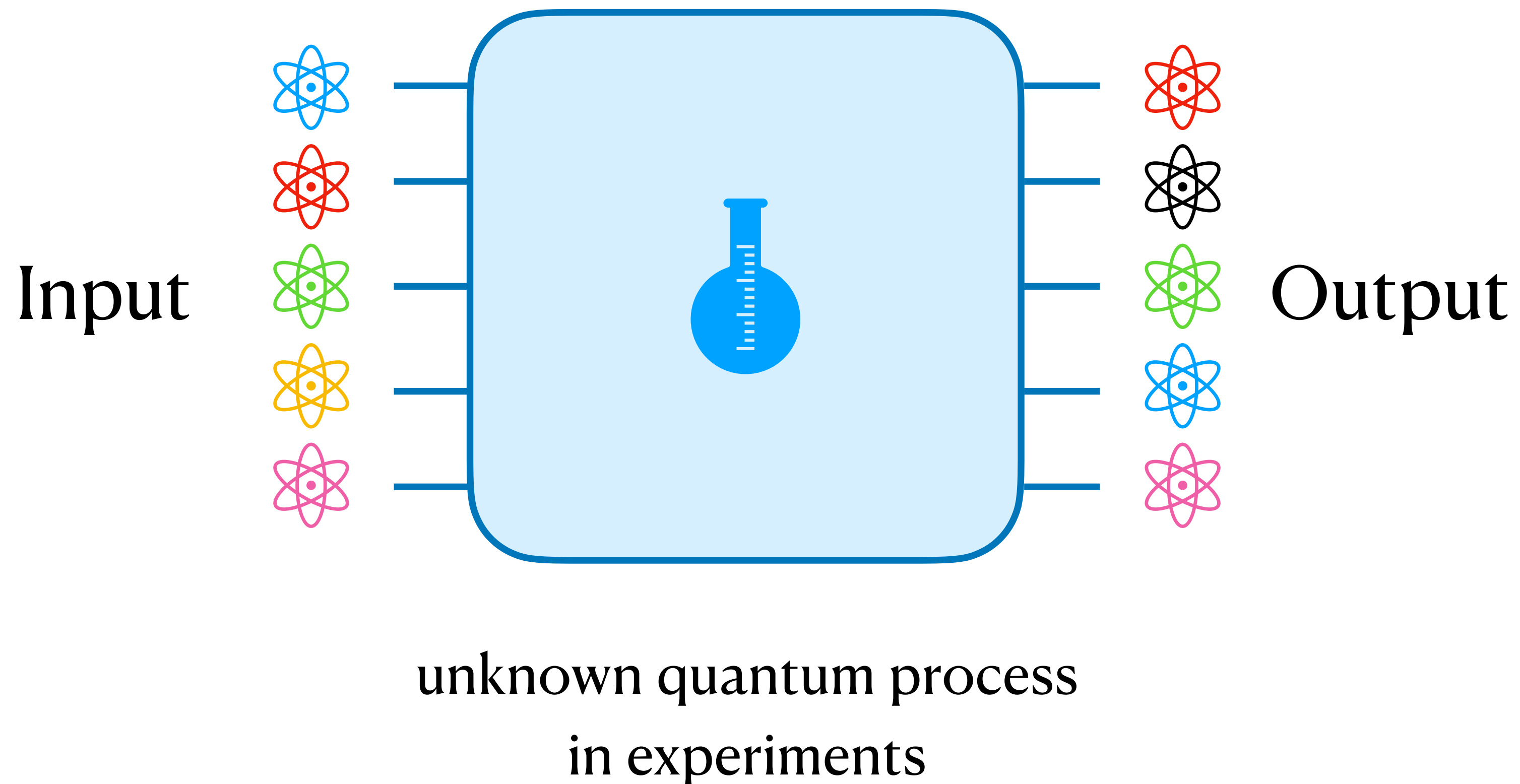


Matthias Caro

Motivation

How to learn from Nature?

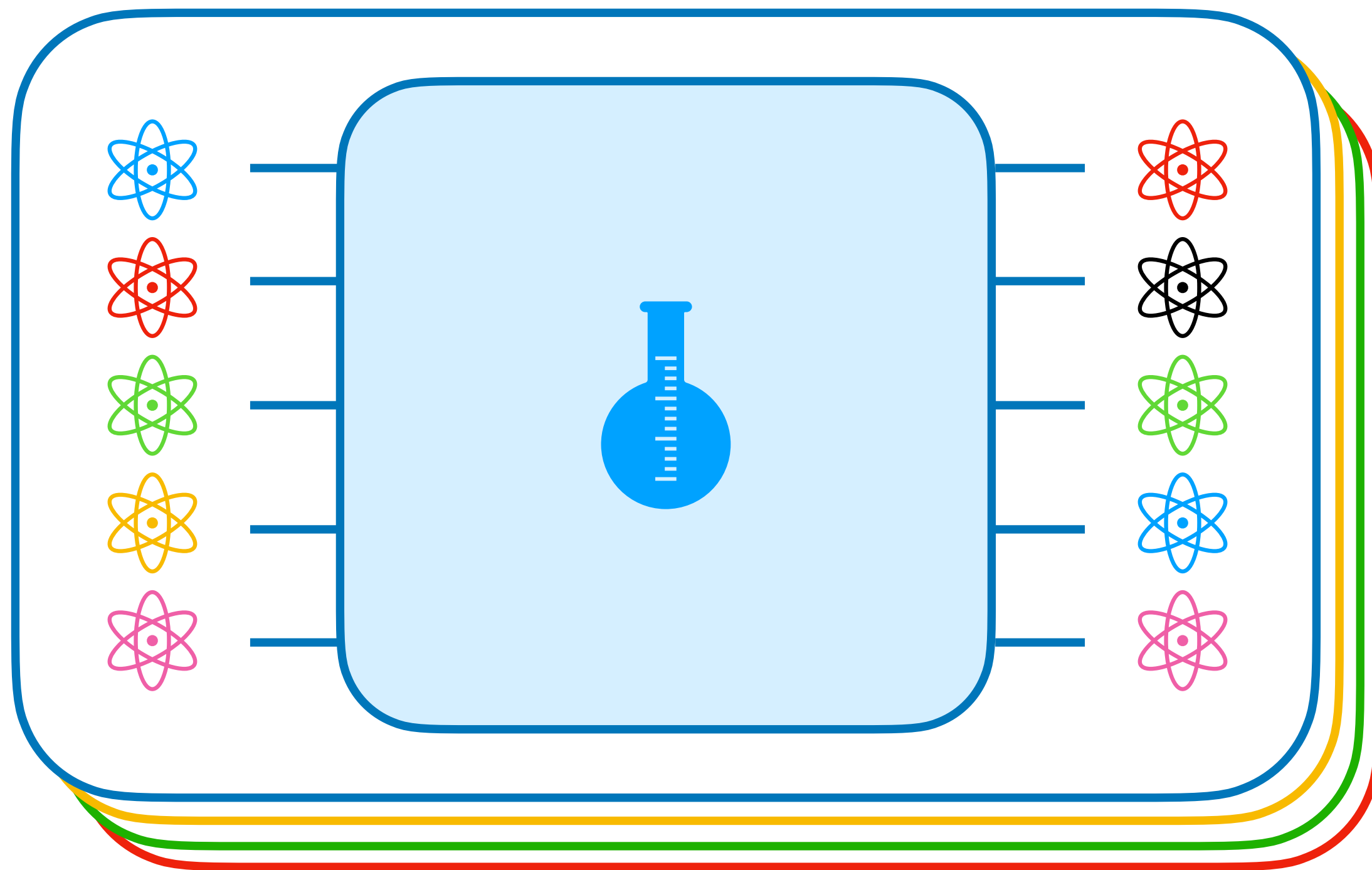
- Do experiments



Motivation

How to learn from Nature?

- Do experiments => Collect many samples

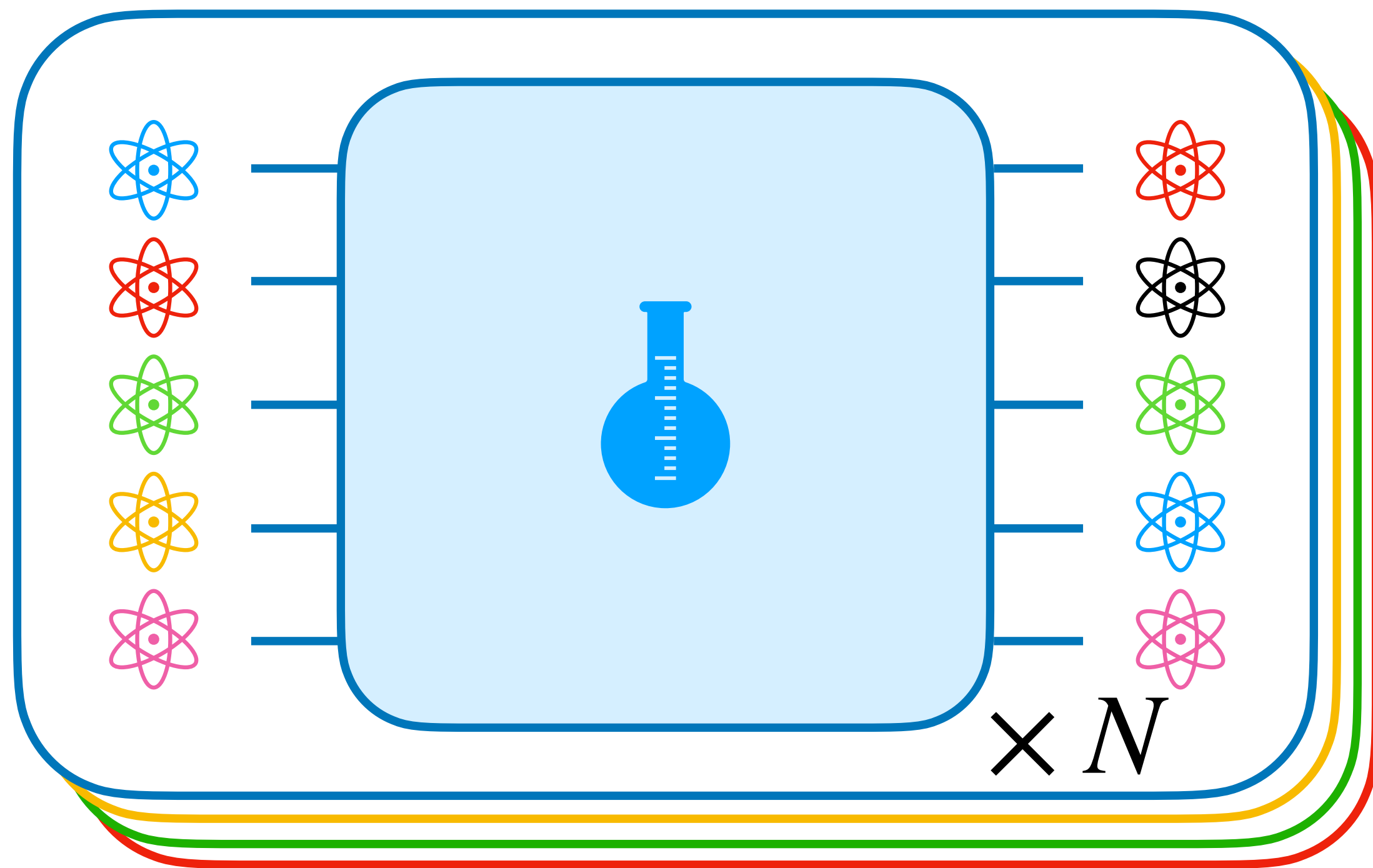


$\times N$

Motivation

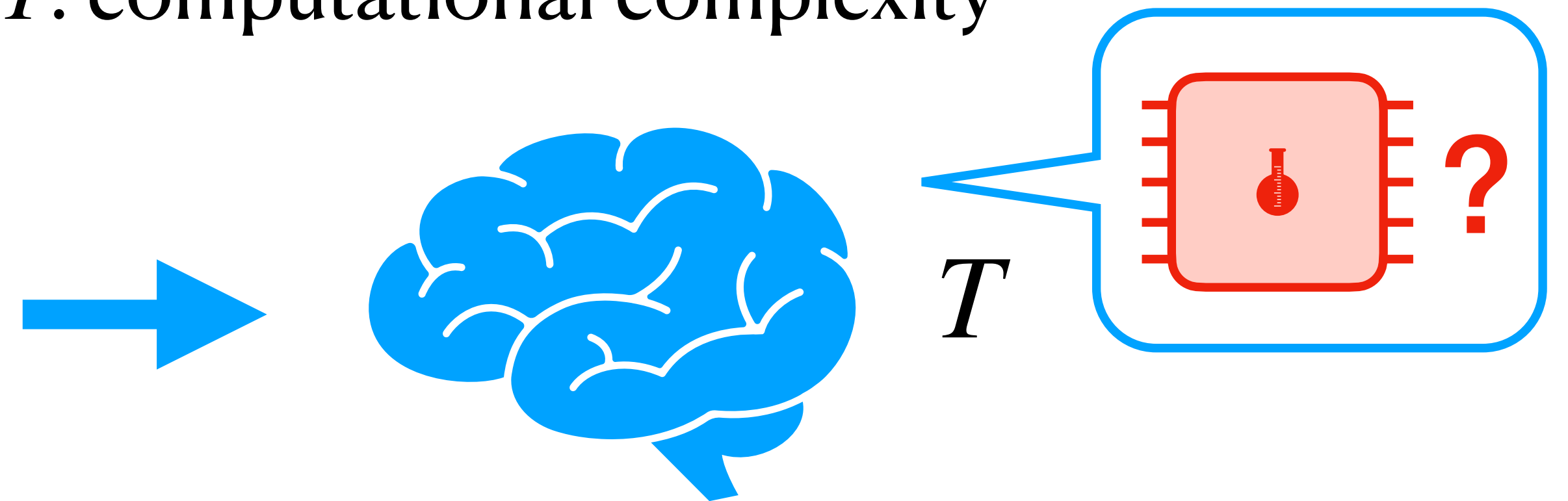
How to learn from Nature?

- Do experiments => Collect many samples => Try to learn the underlying mechanism



N : sample complexity

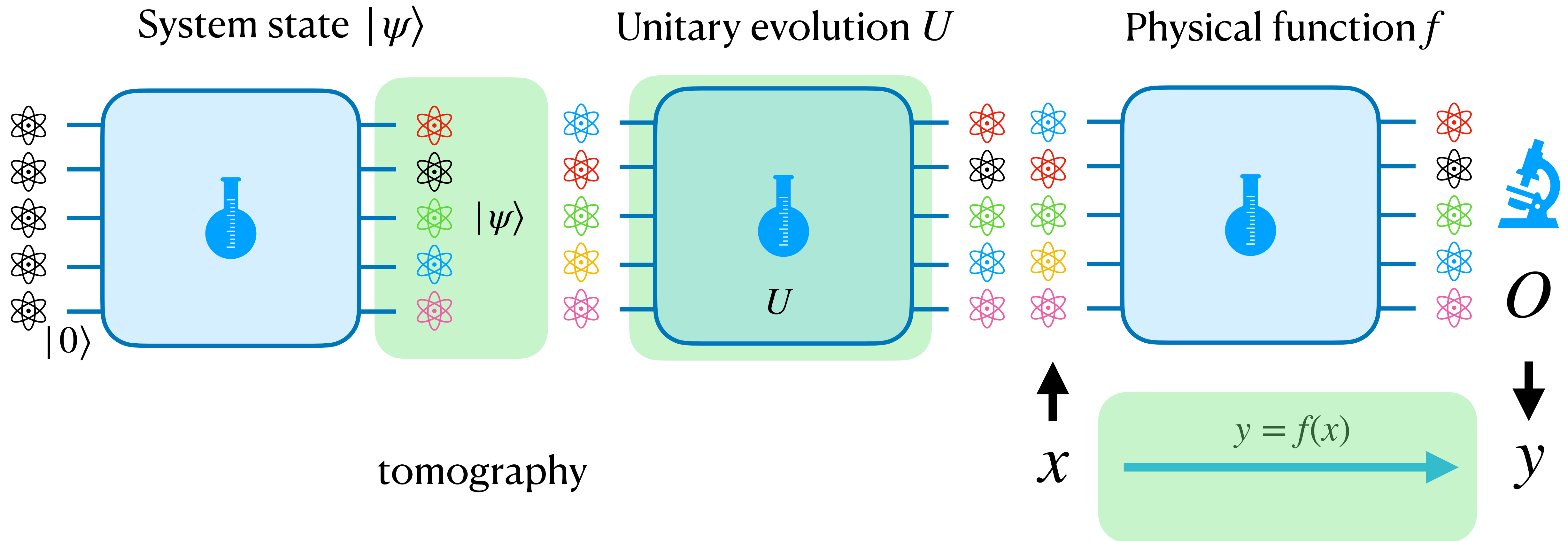
T : computational complexity



$$d(\text{blue chip}, \text{red chip}) \leq \epsilon \text{ w.h.p}$$

Motivation

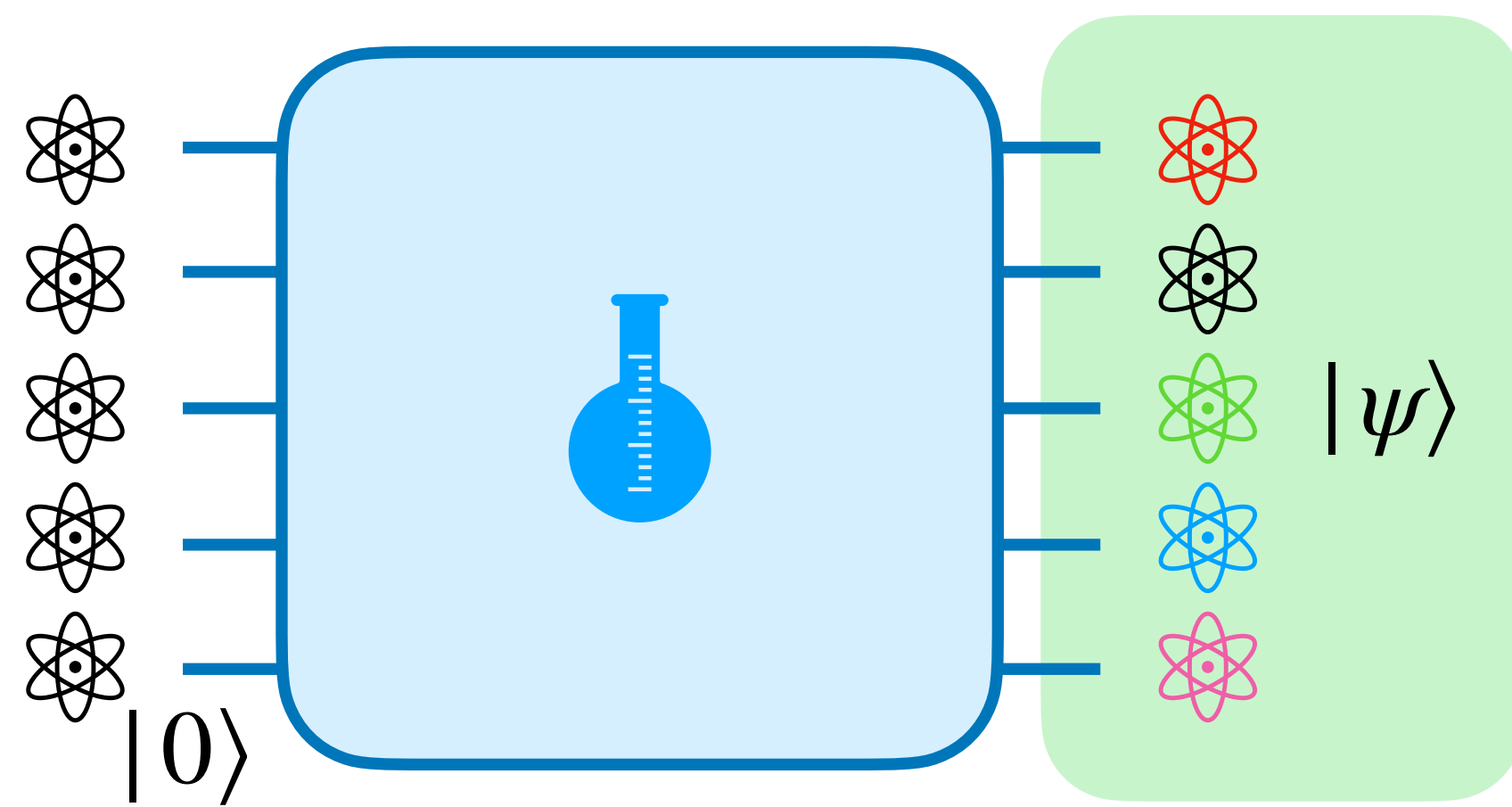
What to learn?



Motivation

Learning is hard in general!

System state $|\psi\rangle$

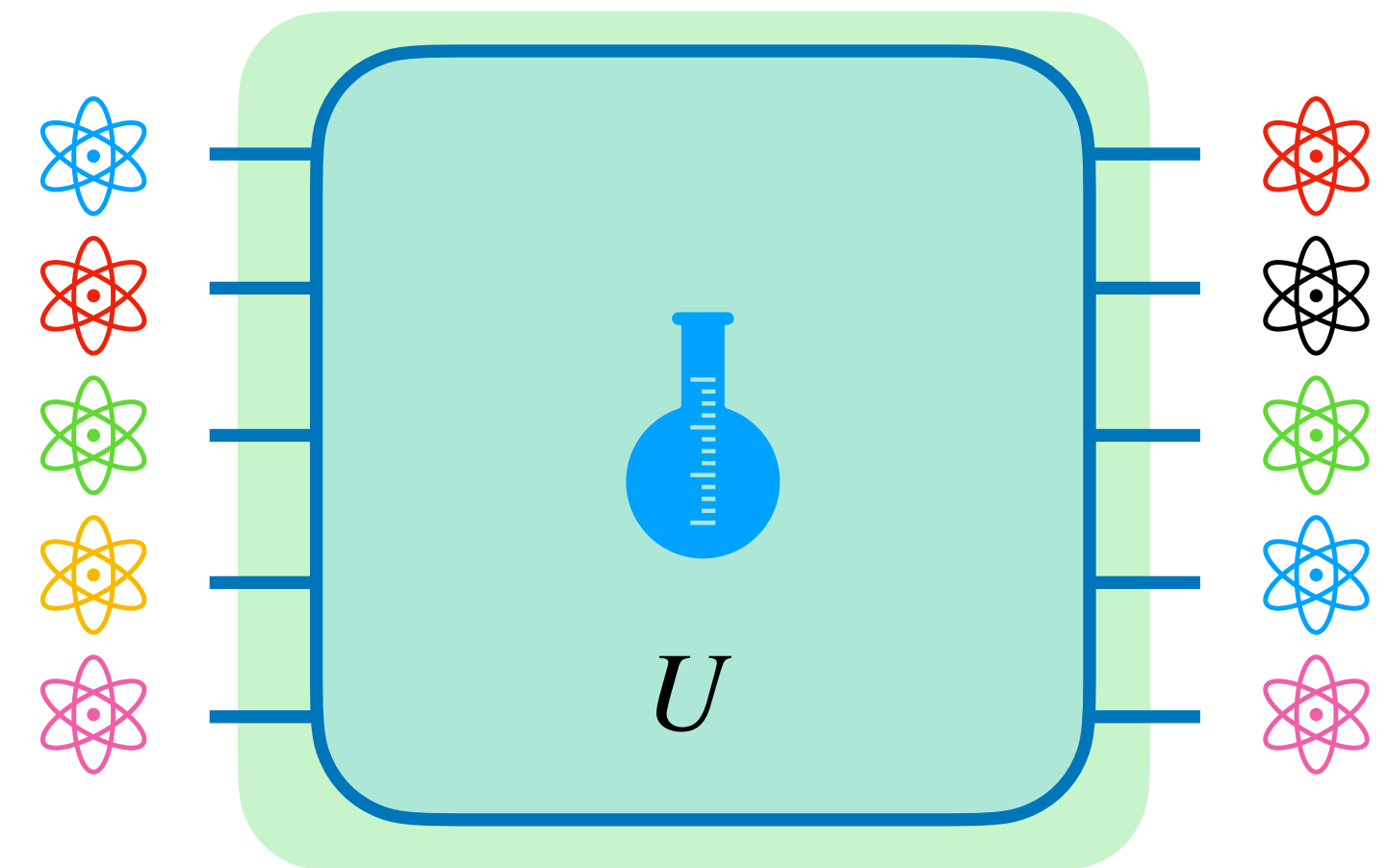


$$N = \Theta(2^n)$$

How is learning even possible?

$$n \sim 10^{23}!$$

Unitary evolution U

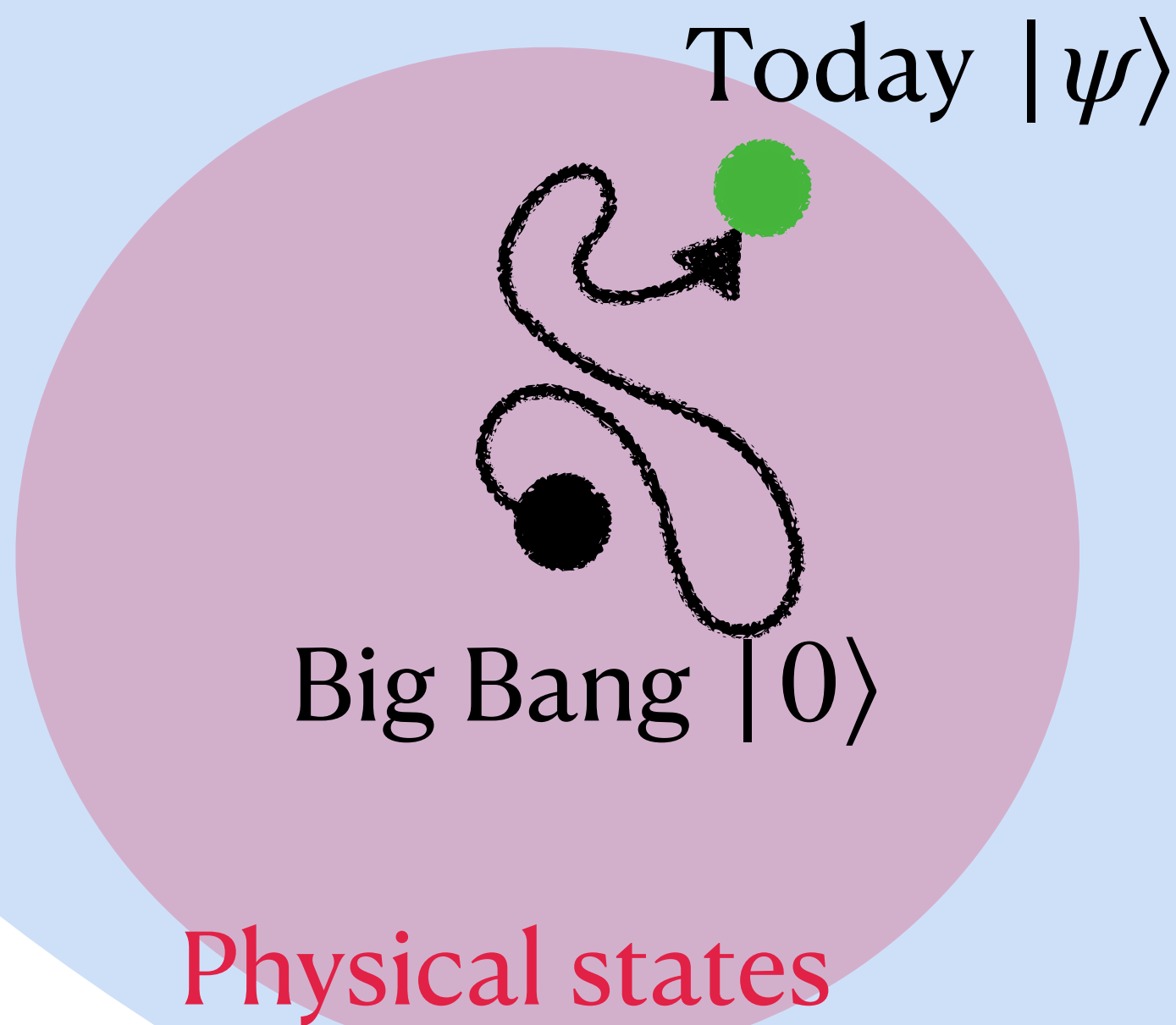


$$N = \Theta(4^n)$$

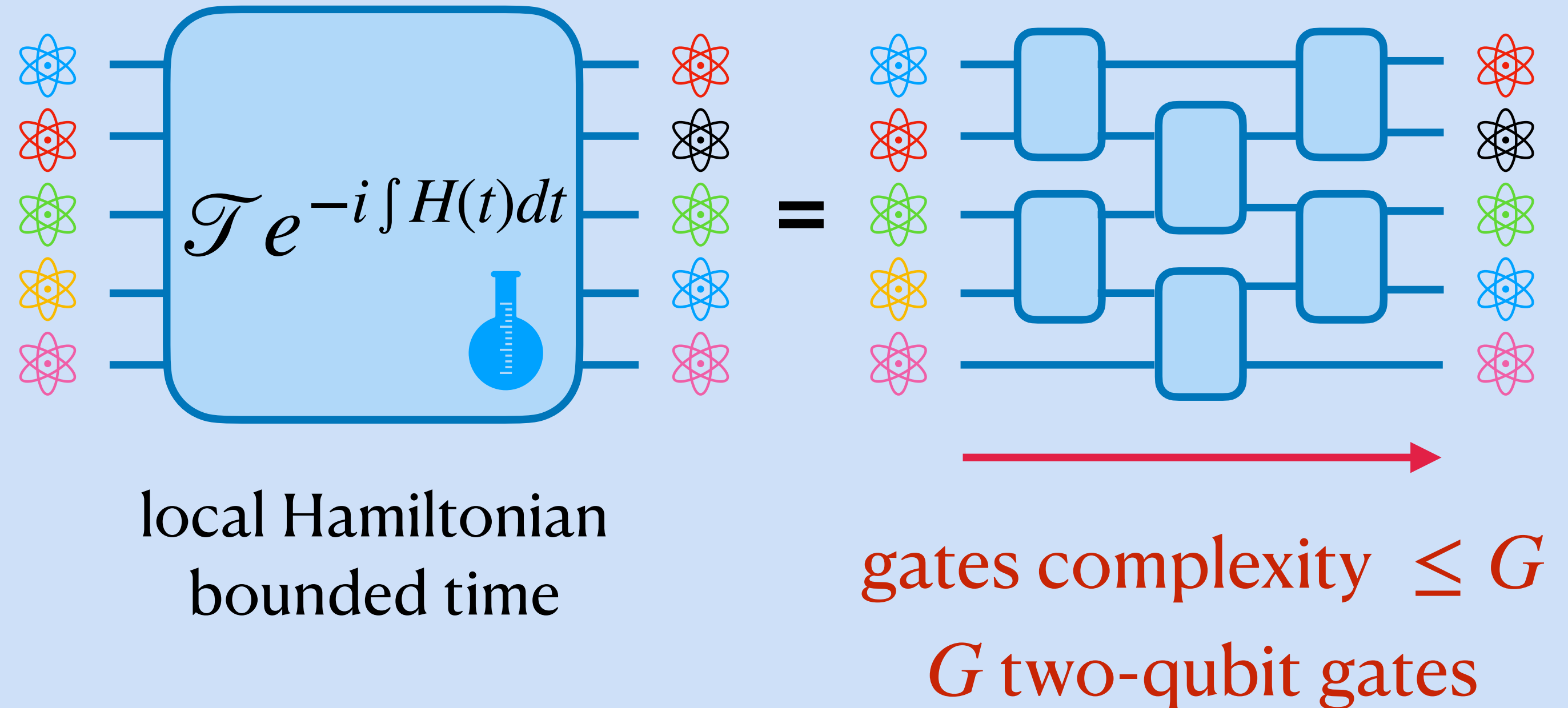
Illusive Hilbert space $\sim 2^n$

Motivation

Physical constraints

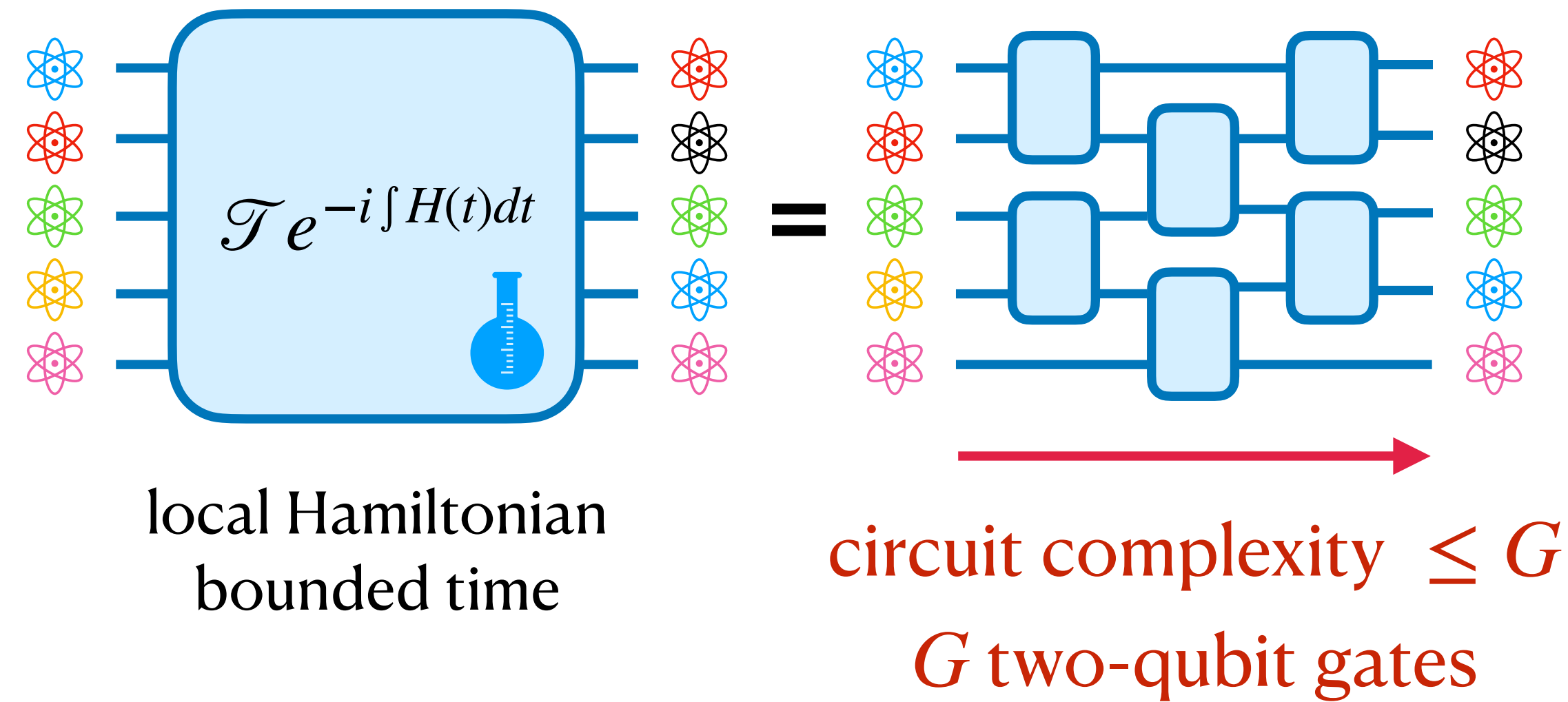


reachable in bounded time



In this work, we don't need "geometric" locality, nor discrete gate-set

Main Question



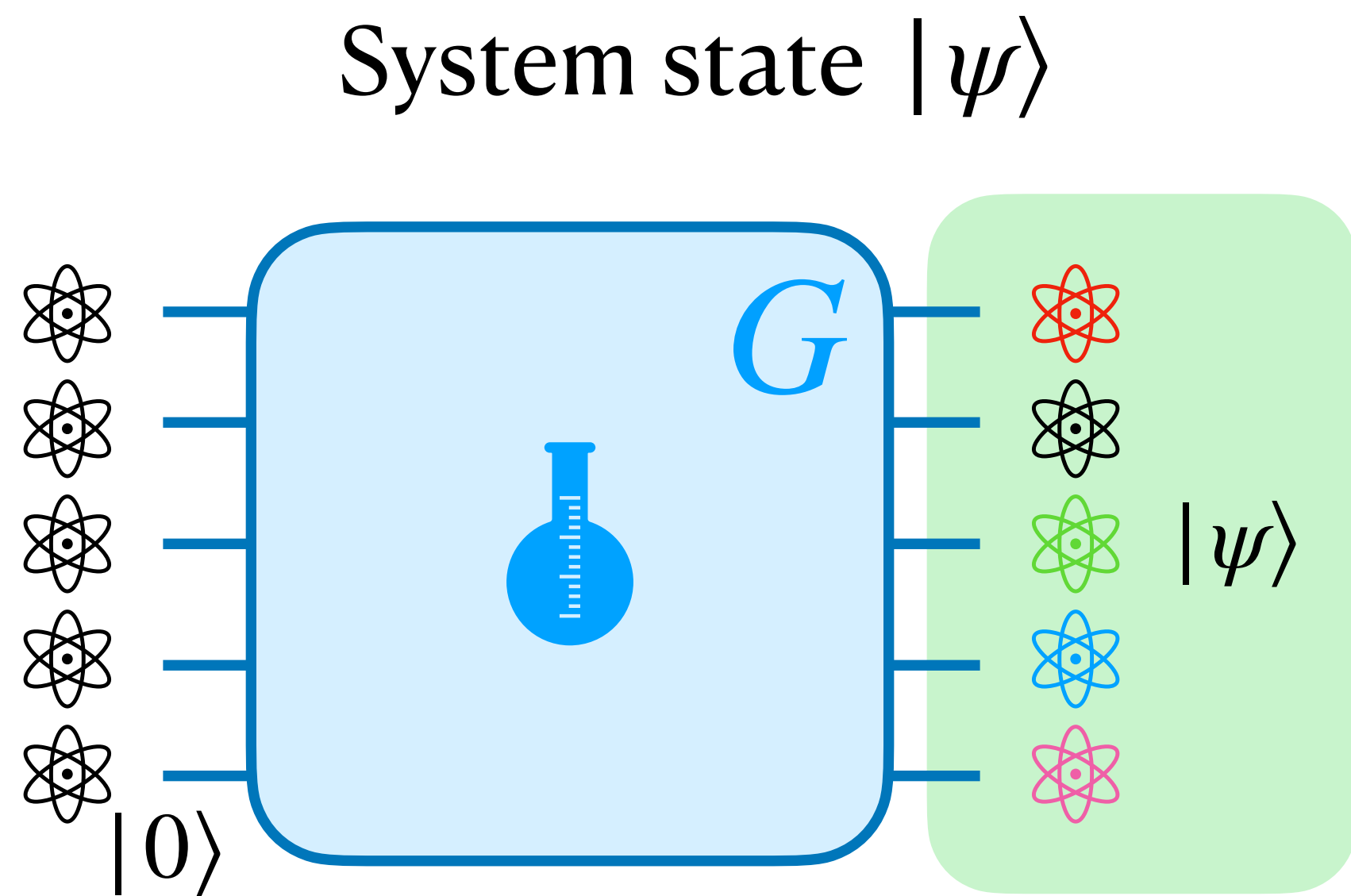
Can we efficiently learn states/unitaries of bounded gate complexity?

Applications on near-term quantum devices: **limited G** (both digital and analog)

Relating different notions of complexity: the complexity of **learning** and **creating**

Results

sample complexity



Learning $|\psi\rangle$ in ϵ trace distance
requires $N = \tilde{\Theta} \left(\frac{G}{\epsilon^2} \right)$ samples.

1. Complexity of learning = complexity of creating (information theoretically)
2. Completely independent of system size n . Can learn $n = 10000, G = 10$ states!
3. Non-adaptive/incoherent scheme is already optimal.

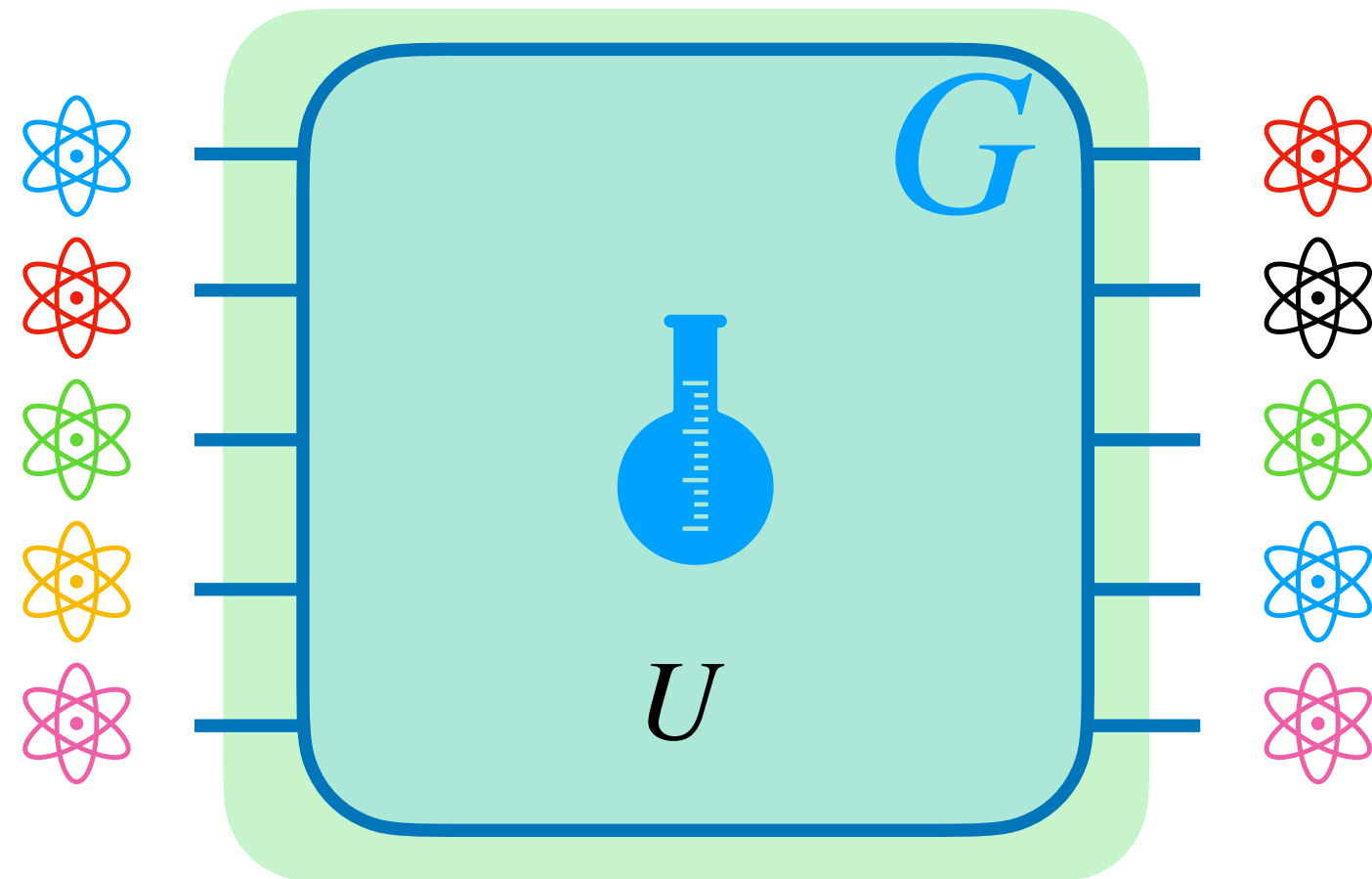
Results

sample complexity

= trace distance between Choi states

$$d_{\text{avg}}(U, V) = \sqrt{\mathbb{E}_{|\psi\rangle} [d_{\text{tr}}(U|\psi\rangle, V|\psi\rangle)^2]}$$

Unitary evolution U

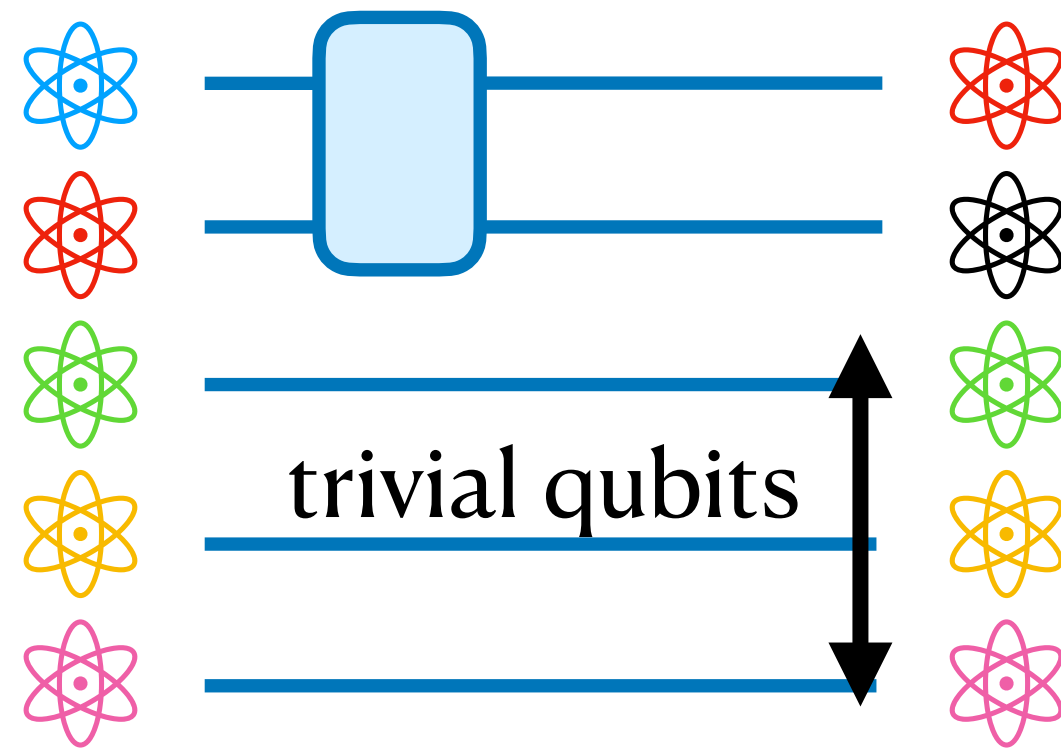


Learning U in average-case distance
requires $N = \tilde{\Theta}(G)$ queries.

1. Non-adaptive/incoherent query is already optimal (in G).
2. Learning in **worst-case distance** (diamond norm) requires $N = \exp(\Omega(\min\{G, n\}))/\epsilon$
Grover
3. ϵ -dependence: $\tilde{O}(\min\{1/\epsilon^2, \sqrt{2^n}/\epsilon\})$, $\Omega(1/\epsilon)$, **Heisenberg scaling** open

Proof sketch

Upper bound: the learning algorithm



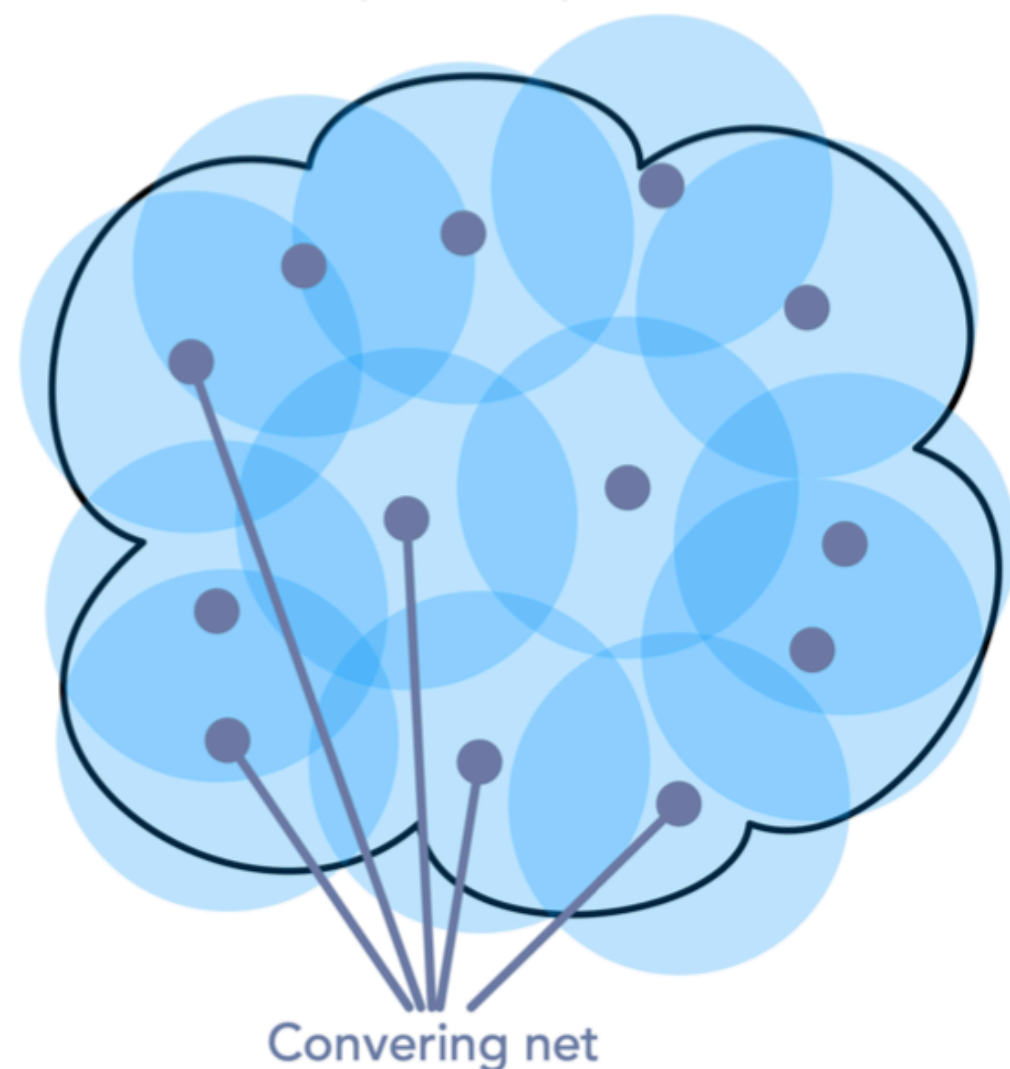
1. **Junta learning:** measure to identify non-trivial qubits, remove n -dependence

2. **Hypothesis selection:** construct a covering net \mathcal{N} that covers the set of G -gate states with ϵ -balls

$$\log |\mathcal{N}| = \tilde{\Theta}(G)$$

3. Find the best candidate by estimating all distances

$$\text{with classical shadow } N = O\left(\frac{\log |\mathcal{N}|}{\epsilon^2}\right) \leq \tilde{O}\left(\frac{G}{\epsilon^2}\right)$$



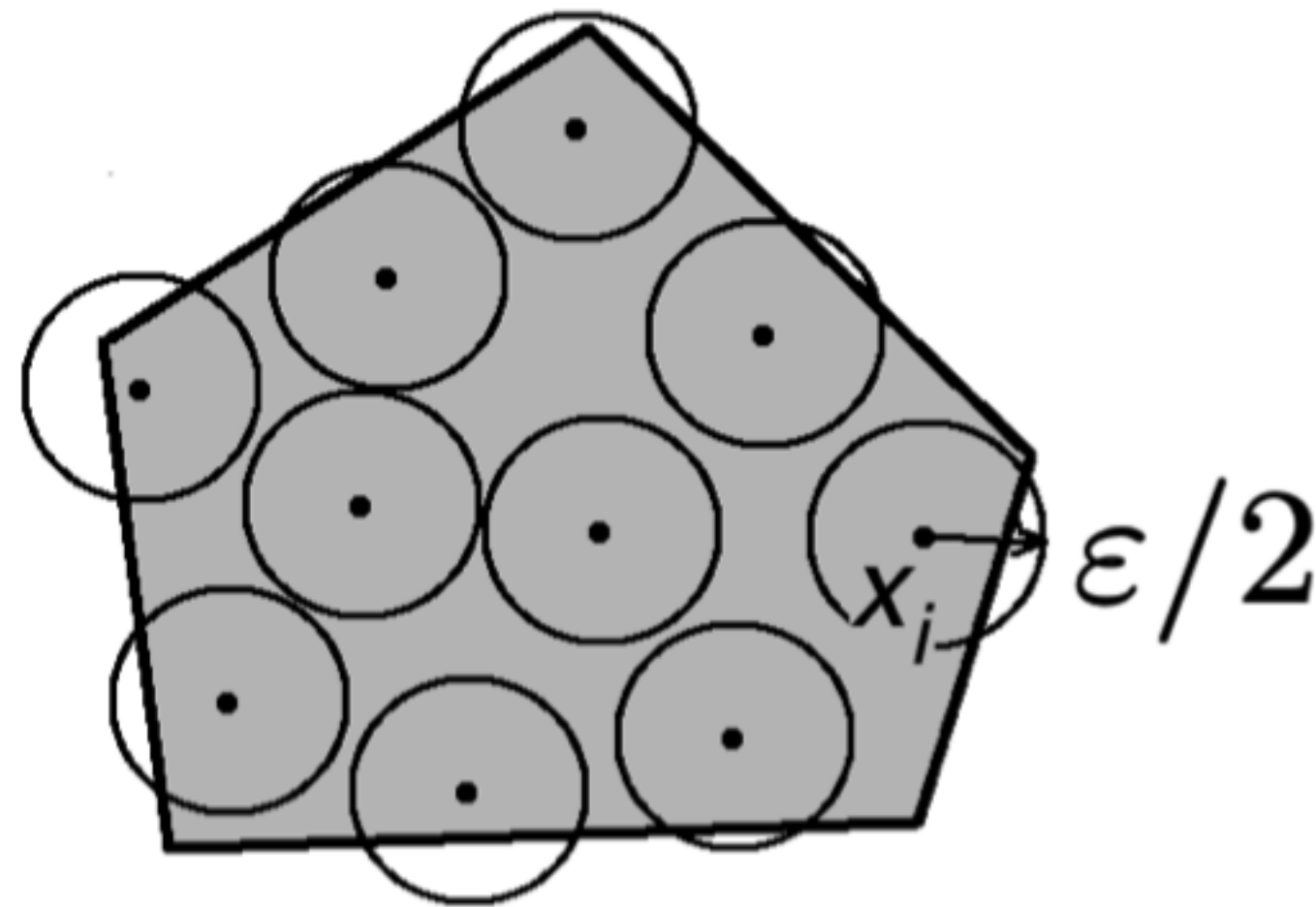
Unitary: $O(G \min\{1/\epsilon^2, \sqrt{2^n}/\epsilon\})$

via Choi states + quantum phase estimation

Chen, Nadimpalli, Yuen, SODA 2023
Bădescu and O'Donnell, STOC 2021
Huang, Kueng, Preskill, Nat Phys 2020

Proof sketch

Lower bound: information theory



1. **Learning \Rightarrow distinguish** elements of a packing net \mathcal{P}

$$\log |\mathcal{P}| = \tilde{\Theta}(G)$$

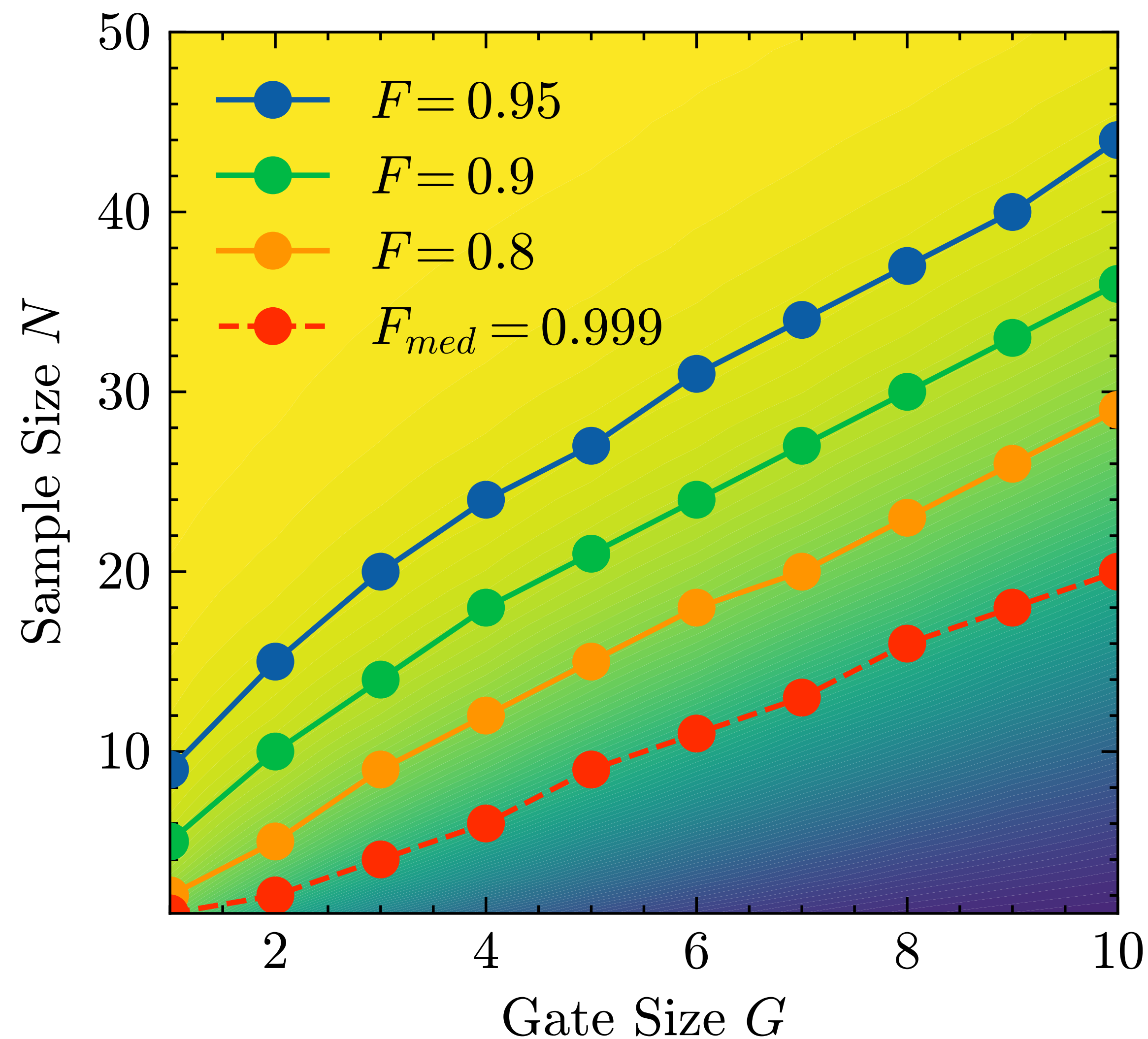
2. **Fano's inequality:** distinguishing requires $\Omega(\log |\mathcal{P}|)$ bits of info

3. **Holevo's theorem:** each sample gives $\tilde{O}(\epsilon^2)$ bits of

$$\text{info} \Rightarrow N \geq \Omega \left(\frac{\log |\mathcal{P}|}{\epsilon^2} \right) = \tilde{\Omega} \left(\frac{G}{\epsilon^2} \right)$$

$\log |\mathcal{P}| \approx \log |\mathcal{N}|$: a general way to prove matching sample complexity bound

Numerical experiments



Learning random G -gate
states on $n = 10000$ qubits

Sample linear in G , runtime e^G

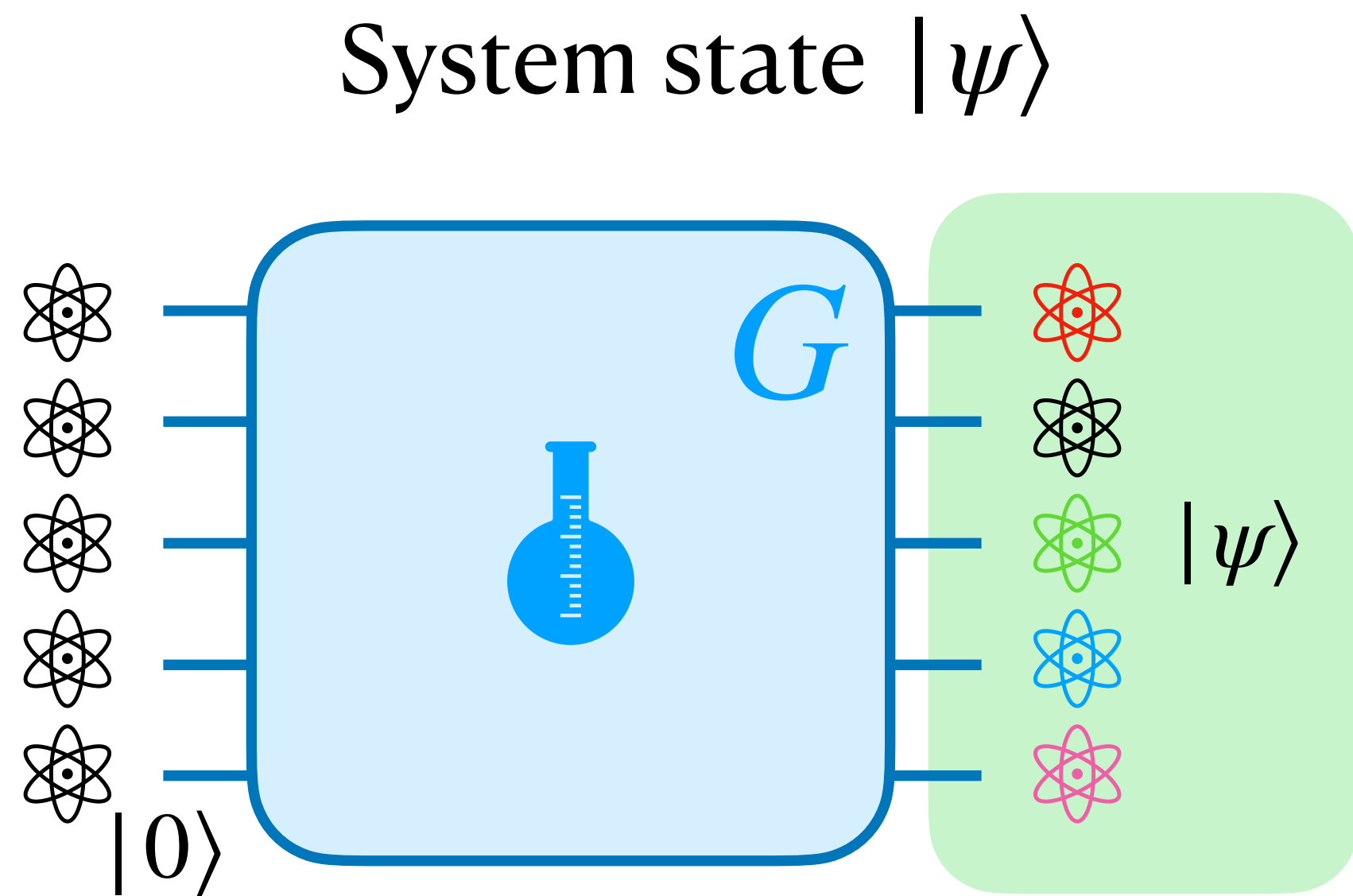


GitHub

haimengzhao/bounded-gate-tomography

Results

computational complexity

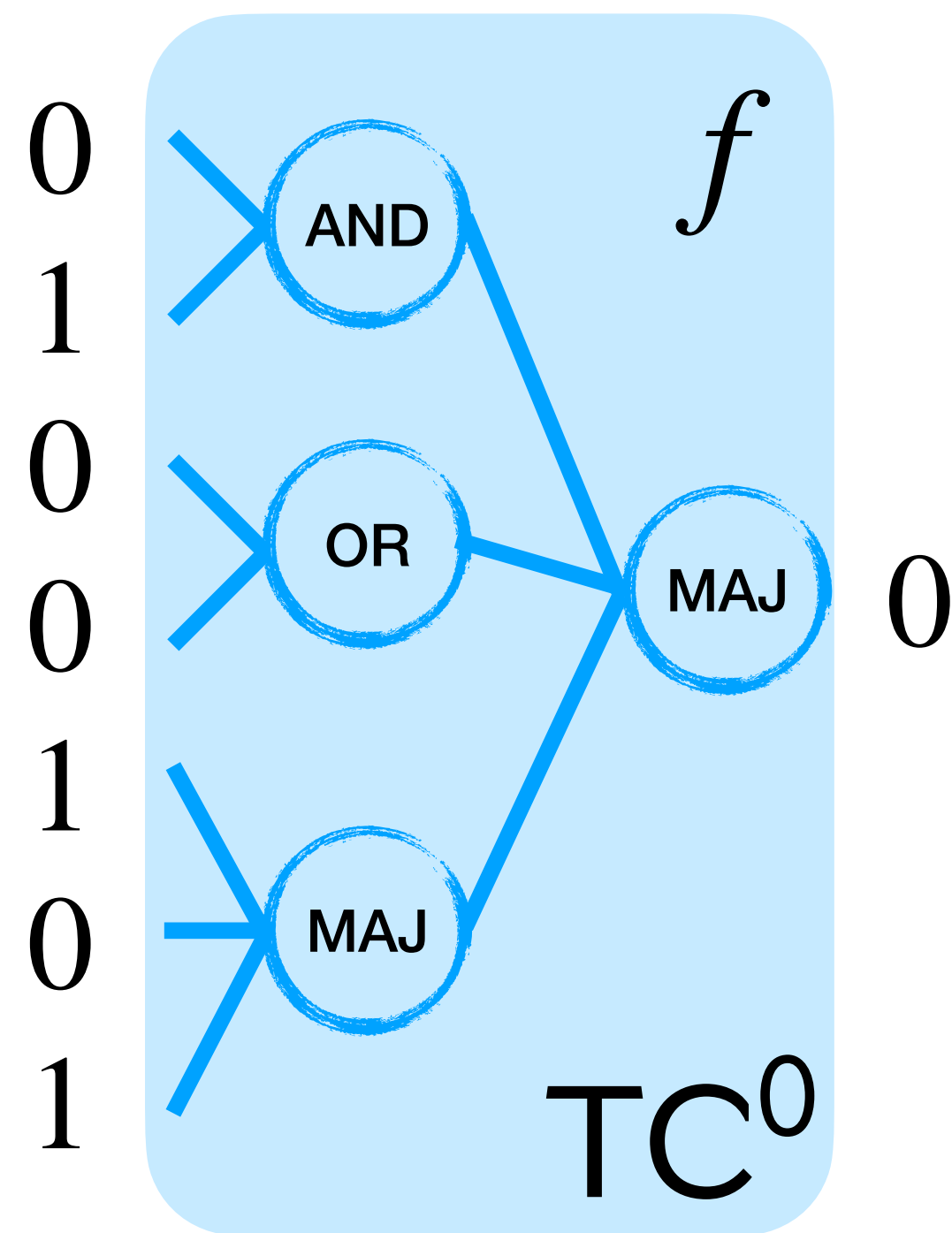


Learning $|\psi\rangle$ in ϵ trace distance
requires $T = \exp(\Omega(\min\{G, n\}))$ time,
if RingLWE is sub-exponential hard.

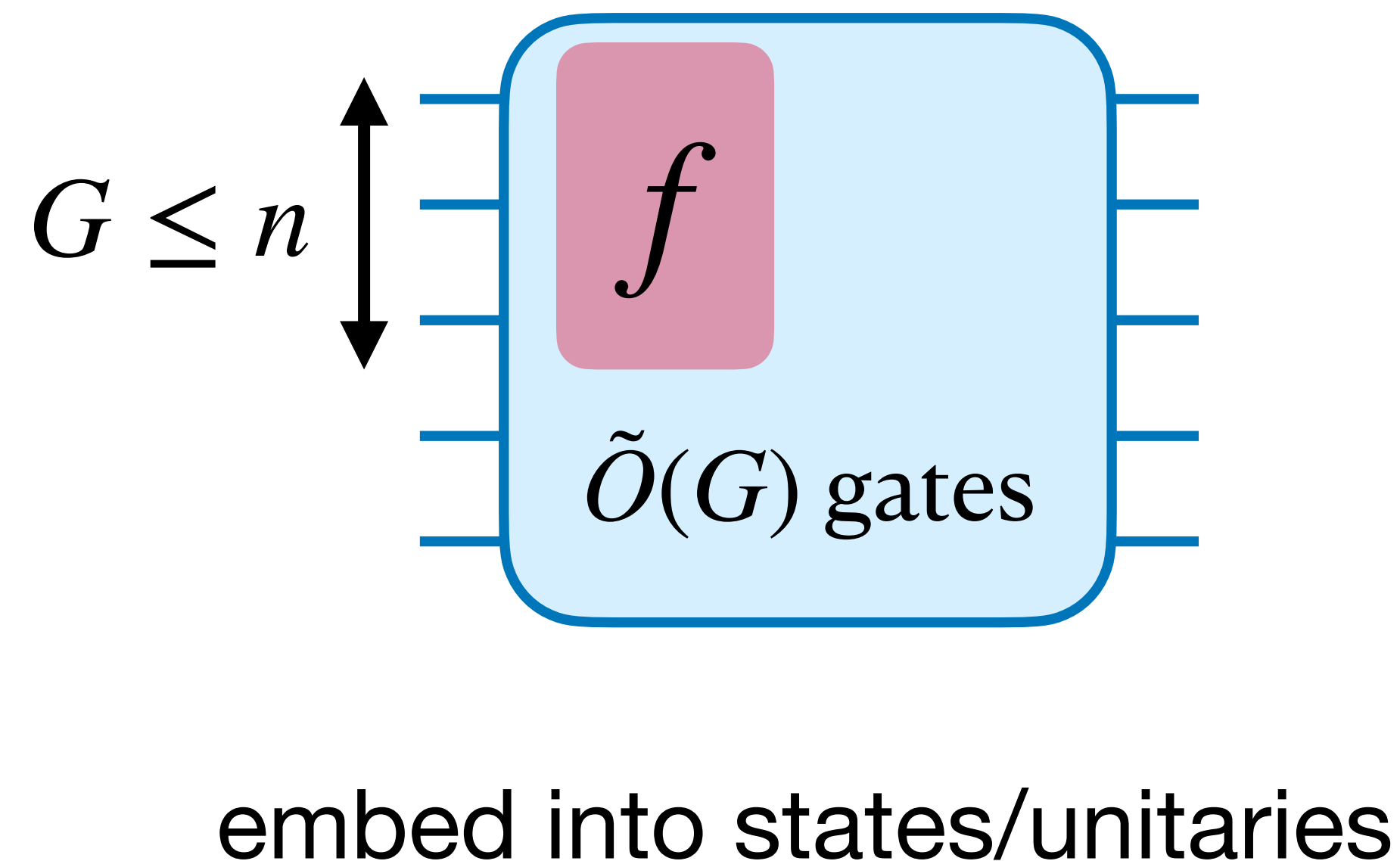
1. Complexity of learning = $e^{\text{complexity of creating}}$ (computationally), efficient $\log n$ gates
2. Even for quantum learners: RingLWE is expected to be hard for quantum computers.
3. Worst-case statement: efficiency possible with additional assumptions.
4. Same for average-case unitary learning

Proof sketch

computational complexity



pseudorandom functions



Learning breaks PRF/PRS

requires $T = e^{\Omega(G)}$ time

$$G \leq n$$

Message



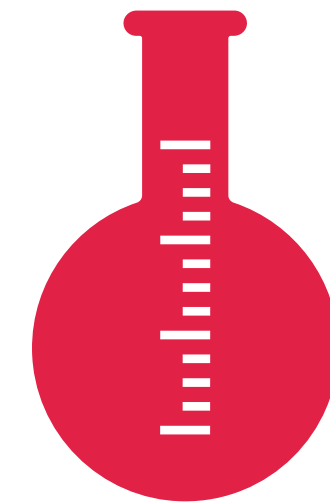
sample size

gate complexity

$$N \approx \log T \approx G \approx t$$

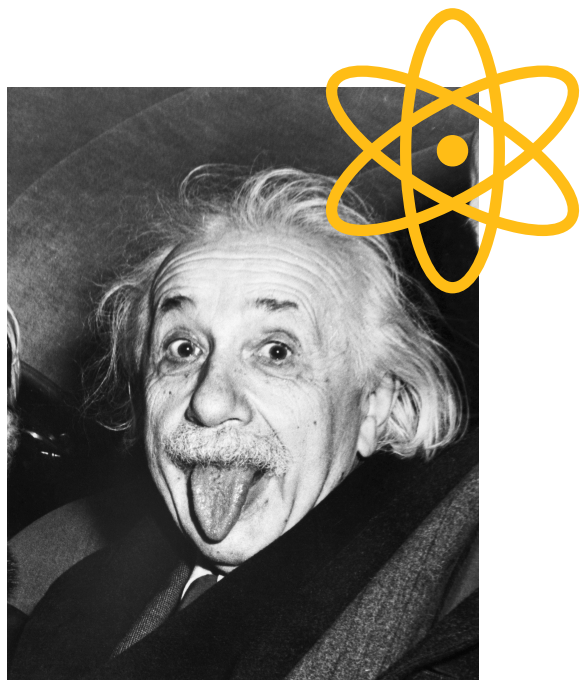
compute time

evolution time



(Brown-Susskind conjecture)

Learning **physical** states/unitaries is **information-theoretically easy**,
but computationally hard!

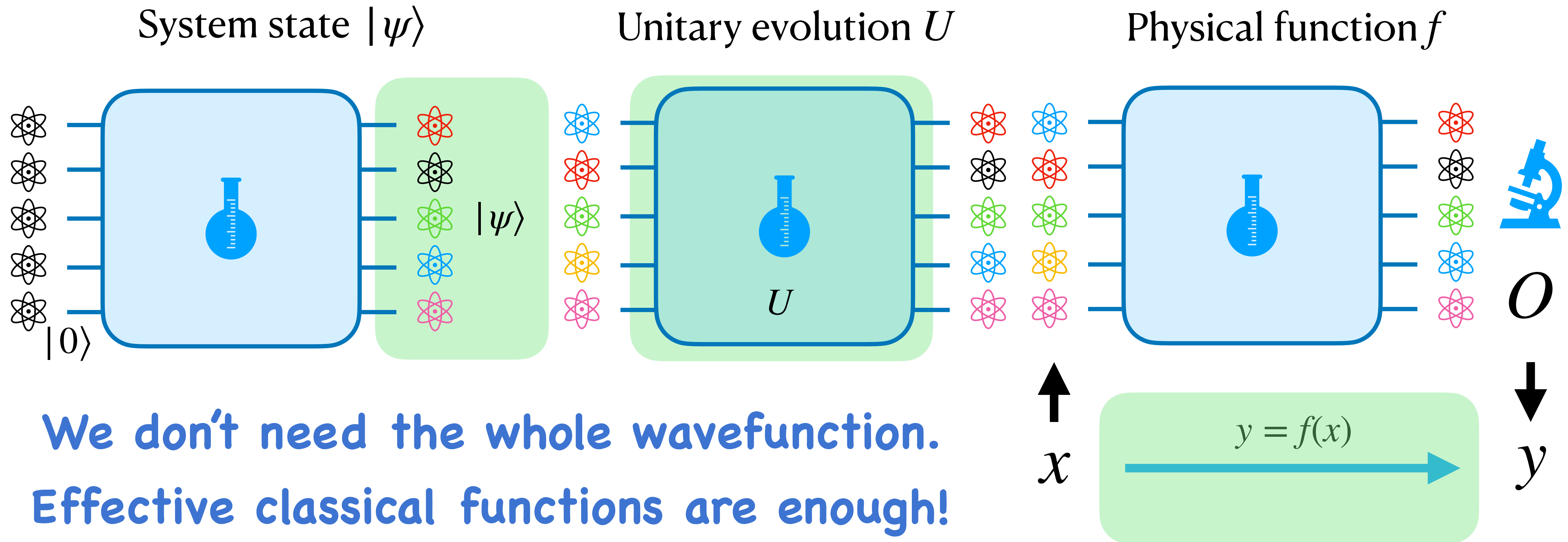


In scientific discovery, **a few samples might already be enough**,
but coming up with a theory requires some real genius!

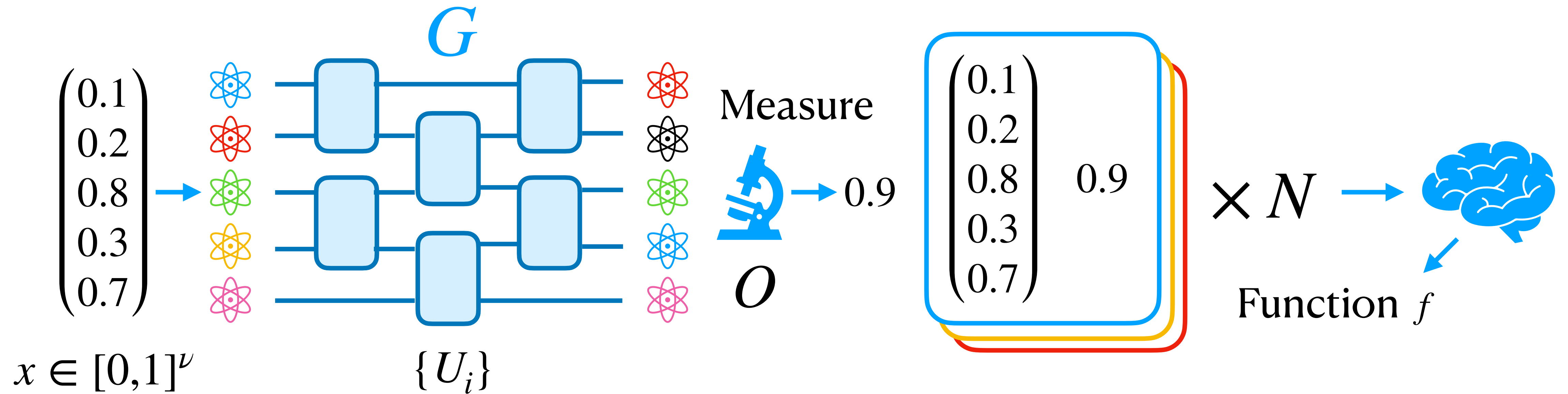
Reminder:
Doing math proofs needs
no data but is NP-hard!

Motivation

What to learn?



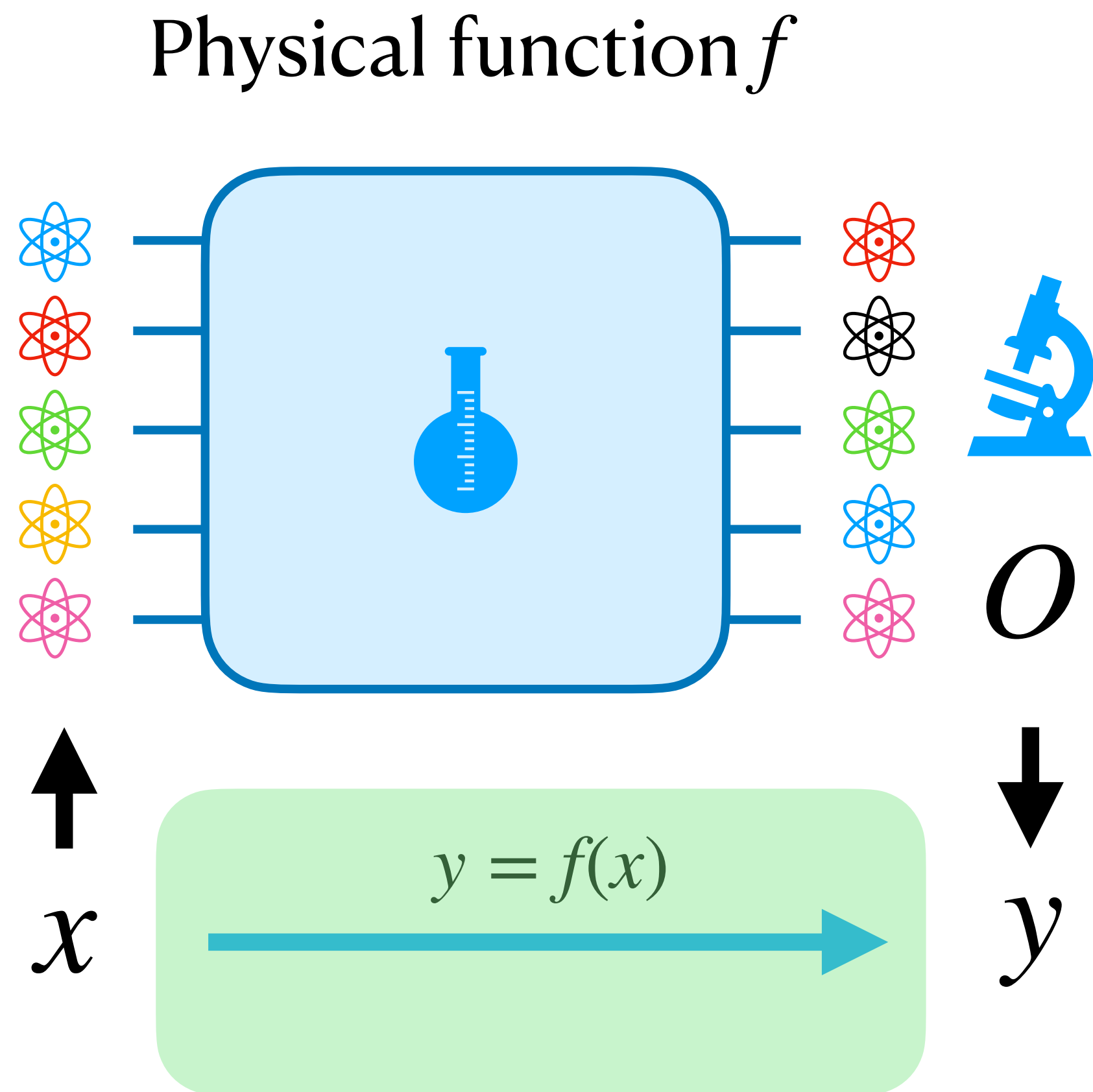
Physical functions



$$f(x) = \langle 0^n | U^\dagger(x, \{U_i\}) O U(x, \{U_i\}) | 0^n \rangle$$

What kind of functions are physically implementable?

Results



To implement/learn 1-Lip 1-bounded functions in ϵ infinite norm,
 $G, N = \tilde{\Omega}(1/\epsilon^\nu)$ gates/samples are needed

1. Certain well-behaved function class is not physical! $G = \Omega(\exp \nu)$
2. For these functions, **no quantum advantage!** $\tilde{\Theta}(1/\epsilon^\nu)$ with classical ReLU NN.
3. More restricted function classes: possible physicality/advantage Fourier integrable $O(1/\epsilon^2)$
Gonon and Jacquier 2023
4. Complement universal approximation theorems of QML Pérez-Salinas et al., PRA 2021
Schuld, Sweke, Meyer, PRA 2021
Manzano, Dechant, Tura, Dunjko 2023

Proof Sketch

$$U_i |\psi\rangle = \begin{pmatrix} U_i^{11} & U_i^{12} \\ U_i^{21} & U_i^{22} \end{pmatrix} \begin{pmatrix} \psi_1 \\ \psi_2 \end{pmatrix} = \begin{pmatrix} U_i^{11}\psi_1 + U_i^{12}\psi_2 \\ U_i^{21}\psi_1 + U_i^{22}\psi_2 \end{pmatrix}$$

$$f(x) = \langle 0^n | U^\dagger(x, \{U_i\}) O U(x, \{U_i\}) | 0^n \rangle$$

$f(x)$ is a **polynomial** of $\{U_i\}$ with degree **at most $2G$**

\Rightarrow fat-shattering/pseudo dimension $\leq \tilde{O}(G)$

Approximating 1-Lip 1-bounded functions require $\Omega(1/\epsilon^\nu)$

$\Rightarrow G, N \geq \tilde{\Omega}(1/\epsilon^\nu)$

Summary and Outlook

- Learning complexity = circuit complexity (information theoretically)
- Learning complexity = $e^{\text{circuit complexity}}$ (computationally)
- Possible efficient learning for more restricted states/unitaries? Clifford+T, MPS, Shallow Circuits, Fermion, Boson, etc.
- What physical properties does learning complexity relate to? Geometry? Phase? Thermalization? Entanglement?
- Worst-case unitary learning requires exponential samples. Du, Hsieh, Tao, arXiv last week
- Initiate the study of physical functions. Learning complexity? Q signal processing?
- Information-theoretic quantum no-free-lunch theorem (not covered)
- Technical: matching bounds, Heisenberg scaling, mixed states/channels, etc.